

Efficient Monadic Streams

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FELDSPAR

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Functional Embedded Language for
Digital Signal Processing and Parallelism

FELDSPAR

- ▶ Embedded in Haskell
- ▶ Generates C
- ▶ Aims to raise the level of abstraction for the programming

- ▶ Streams for signal processing

```
data Stream a where
```

```
  Stream :: (s -> (a,s)) -> s -> Stream a
```

Compositional list-like interface

```
mapStream :: (a -> b) -> Stream a -> Stream b
mapStream f (Stream next i) = Stream next' i
  where next' s = case next s of
                    (a,s') -> (f a,s')
```

```
zip :: Stream a -> Stream b -> Stream (a,b)
zip (Stream next1 i1) (Stream next2 i2) = Stream next (i1,i2)
  where next (s1,s2) = case (next1 s1,next2 s2) of
                        ((a,s1'),(b,s2')) -> ((a,b),(s1',s'2))
```

- ▶ The functional stream representation supports *fusion*
- ▶ Simply unfold function definitions
- ▶ Allows for compositional programming without performance loss
- ▶ This transformation happens as a side-effect of how Feldspar is embedded in Haskell.
No extra implementation effort needed.

```
mapStream f (mapStream g (Stream next init))
```

```
==> { unfolding mapStream }
```

```
mapStream f (Stream next' init)
```

```
  where next' s = case next s of (a,s') -> (g a,s')
```

```
==> { unfolding mapStream }
```

```
Stream next'' init
```

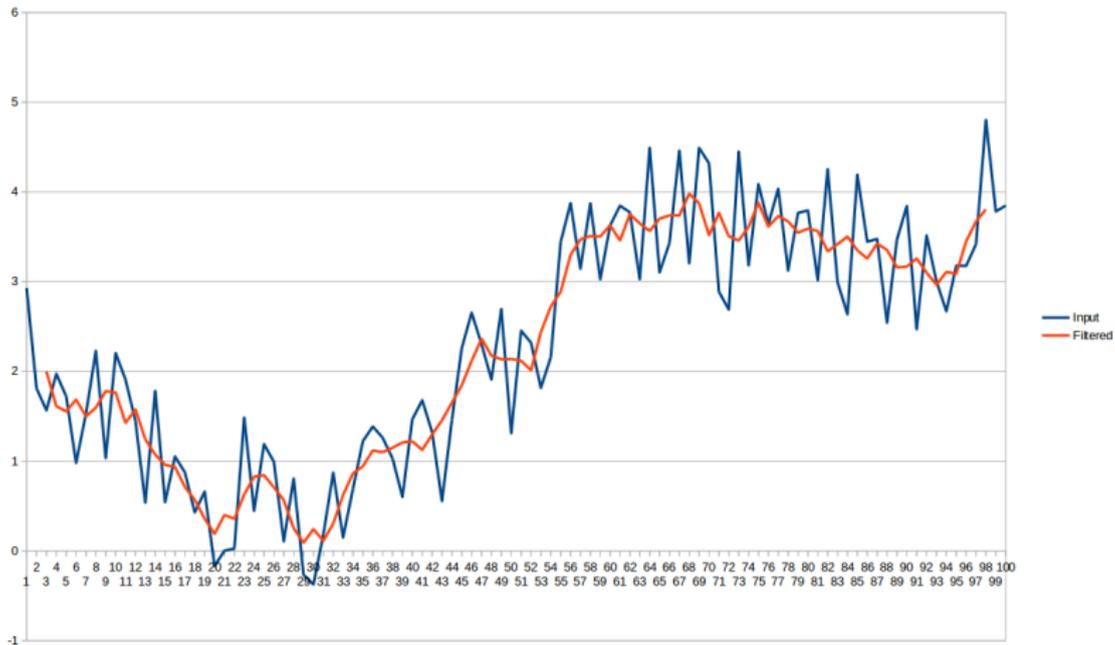
```
  where next' s = case next s of (a,s') -> (g a,s')
```

```
        next'' s = case next' s of (b,s') -> (f b,s')
```

```
==> { unfolding next' and evaluation }
```

```
Stream next'' init
```

```
  where next'' s = case next s of (a,s') -> (f (g a),s')
```



```
movingAvg :: Fractional a => Data Int -> Stream a -> Stream a
movingAvg n (Stream step init) = Stream step' init'
  where
    init' = (init, replicate n 0)
    step' (s, window) =
      let (a, s') = step s
          window' = init window ++ singleton a
      in (avg window, (s', window'))
    avg w = sum (elems w) / fromIntegral n
```

```

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    avg w = sum (elems w) / fromIntegral n

```

- ▶ The sliding window copied each iteration
- ▶ Costly when window grow large

```
1  for (uint32_t v47 = 0; v47 < 32; v47 += 1)
2  {
3      //Code for computing the average of the window elided
4      for (uint32_t v89 = 0; v89 < v184; v89 += 1)
5      {
6          ((v46).member2).member2[(v63 + 1)] =
7              ((v48).member2).member2[v63];
8      }
9      //Code updating the struct which holds the window elided
10 }
```

```
1  for (uint32_t v47 = 0; v47 < 32; v47 += 1)
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3      //Code for computing the average of the window elided
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5      {
6          ((v46).member2).member2[(v63 + 1)] =
7              ((v48).member2).member2[v63];
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```

- ▶ Line 4-8 is a loop which copies the window

We would like a representation where we can avoid copying the buffer.

Monads to the rescue

```
data Stream a = Stream (M (M a))
```

```
data Stream a = Stream (M (M a))
```

- ▶ The monad `M` in `Feldspar` is similar to `IO` in Haskell
- ▶ Provides mutable references and arrays

```
1 enumFrom :: Data Int -> Stream Int
2 enumFrom n = Stream $ do
3   r <- newRef n
4   loop $ do
5     i <- readRef r
6     writeRef r (i+1)
7     return i
```

```
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```

- ▶ Line 3: Initialization
- ▶ Lines 5-7: Loop body
- ▶ loop = return

```
data Stream a = Stream (M (M a))
```

- ▶ The outer monad performs initialization
- ▶ The inner monad is the loop body.

```
mapStream :: (a -> b) -> Stream a -> Stream b
mapStream f (Stream init) = Stream $ do
  body <- init
  loop $ do
    a <- body
    return (f a)

zip :: Stream a -> Stream b -> Stream (a,b)
zip (Stream init1) (Stream init2) = Stream $ do
  next1 <- init1
  next2 <- init2
  loop $ do
    a <- next1
    b <- next2
    return (a,b)
```

Fusion

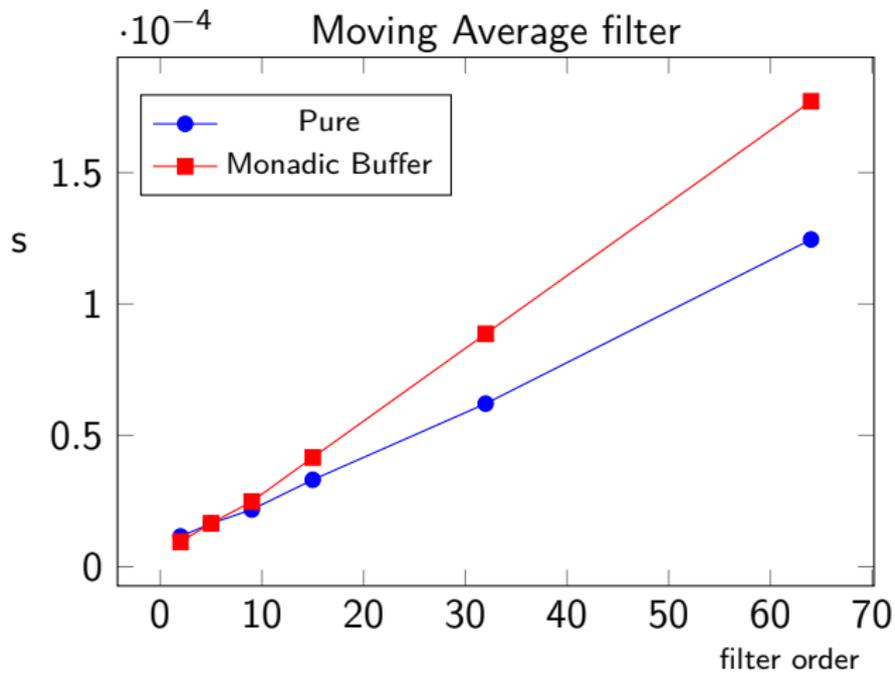
- ▶ The new stream type also supports fusion
 - ▶ Unfolding of definitions, just like in the functional representation
 - ▶ Two of the monad laws: associativity and left identity
- ▶ In Feldspar, the two monad laws are also applied as a side-effect of the way it is embedded in Haskell

```
1 movingAvg :: Int -> Stream Double -> Stream Double
2 movingAvg n s =
3   recurrence (listArray (0,n-1) (replicate n 0.0)) s
4             (\input -> sum (elems input) / fromIntegral n)
```

```
1 recurrence :: Array Int a -> Stream a ->
2             (Array Int a -> b) ->
3             Stream b
4 recurrence ii (Stream init) mkExpr = Stream $ do
5     next <- init
6     ibuf <- initBuffer ii
7     loop $ do
8         a <- next
9         putBuf ibuf a
10        b <- withBuf ibuf $ \ib ->
11            mkExpr ib
12        return b
```

```
1  v14 = 0;
2  copy(v7, zeros);
3  for (uint32_t v24 = 0; v24 < 32; v24 += 1) {
4      v48 = v0[v25];
5      v27 = v14;
6      v14 = ((v27 + 1) % 8);
7      v7[v27] = v48;
8      // Code computing the average elided
9  }
```

- ▶ No copying
- ▶ Line 7 updates the window



- ▶ The cyclic buffer uses **modulus** for each element
- ▶ More expensive than copying

Making a faster filter

Although the mutable buffer was slower, mutable streams allow for many variations.

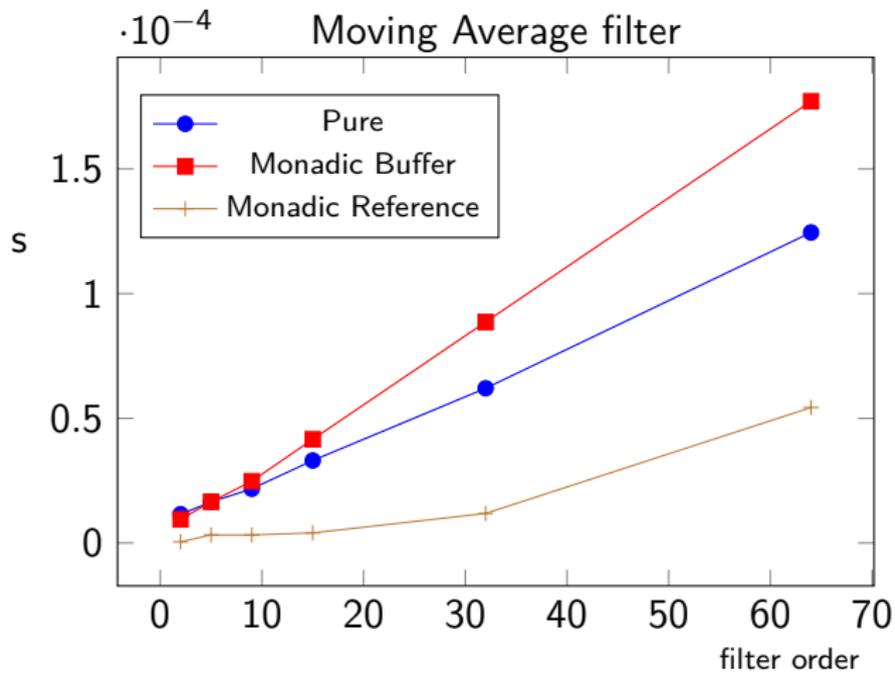
For instance:

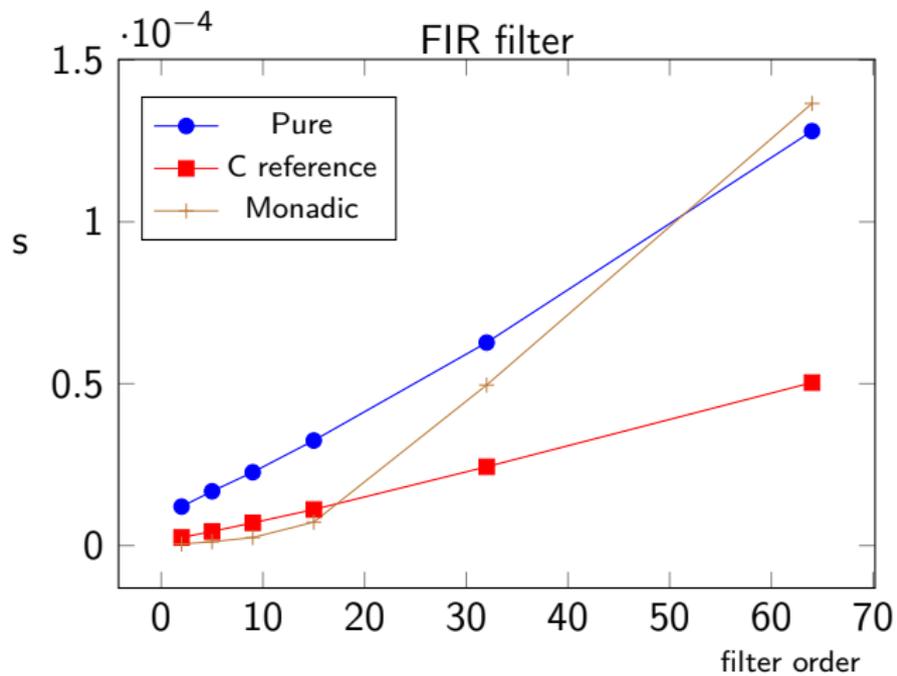
- ▶ Unroll the buffer to keep each element in a reference
- ▶ Effectively means that each element is kept in a register

```

recurrenceS :: (Type a, Type b) =>
  [Data a] -> Stream (Data a) ->
  ([Data a] -> Data b) ->
  Stream (Data b)
recurrenceS ii (Stream init) mkExpr = Stream $ do
  next <- init
  ris <- mapM newRef ii
  loop $ do
    a <- next
    if (not $ null ii) then pBuf ris a else return ()
    b <- wBuf ris $ \ib ->
      return $ mkExpr ib
  return b
where
  pBuf rs a = zipWithM (\r1 r2 -> getRef r1 >=> setRef r2)
    (tail $ reverse rs) (reverse rs)
    >> setRef (head rs) a
  wBuf rs f = mapM getRef rs >>= f

```





Conclusion

- ▶ A simple, versatile stream representation
- ▶ Enables mutation for efficiency
- ▶ A declarative, compositional interface
- ▶ Well suited for embedded languages
- ▶ Fusion can happen for free with the right embedding

Unexpectedly

- ▶ Mutation is not always faster

In the paper

- ▶ Avoiding multiple loop variables